



## CAD Surfaces & solid modeling





# **LIEGE** CAD & Computational Geometry CAD Surfaces & solid modeling

- Parametric surfaces
  - (Coons patches) "like" splines for curves
  - Tensor product surfaces
- Procedural (non parametric) surfaces
  - Subdivision surfaces
- Solid modeling





#### **CAD Surfaces**

# "Tensor product" surfaces





# Parametric surfaces as a polar form

 $S(u, v) = \sum_{k} N_{k}(u, v) P_{k}$ 

- One shape function per control point
- If  $N_k$  is separable : « Tensor product » surface
  - Combination of elementary curves/shape functions independently defined on *u* and *v*.  $S(u, v) = \sum_{i} \sum_{j} G_{i}(u) H_{j}(v) P_{ij}$

Usually built upon Bézier and B-Splines curves/SFs

• The "unique" shape function is  $N_k(u, v) = G_i(u)H_j(v)$ 

# LIÈGE CAD & Computational Geometry B-Spline surfaces



- B-Splines surfaces uses 1D B-Spline shape fns.
  - Definition as tensor product :

$$S(u, v) = \sum_{i=0}^{n} \sum_{j=0}^{m} N_{i}^{p}(u) N_{j}^{q}(v) P_{ij}$$

Every variable u and v has a degree (p and q) and a nodal sequence U and V:

$$U = \{\underbrace{0, \cdots, 0}_{p+1}, u_{p+1}, \cdots, u_{r-p-1}, \underbrace{1, \cdots, 1}_{p+1}\} \quad (r+1 \text{ nodes})$$
$$V = \{\underbrace{0, \cdots, 0}_{q+1}, v_{q+1}, \cdots, v_{s-q-1}, \underbrace{1, \cdots, 1}_{q+1}\} \quad (s+1 \text{ nodes})$$

- The control points forms a regular net P<sub>ij</sub> (n+1 times m+1) values.
- We have the following relations : r=n+p+1 s=m+q+1 5





## Example of basis functions

 $U = \{0, 0, 0, 0, 1/4, 1/2, 3/4, 1, 1, 1, 1\} \quad p = 3$  $V = \{0, 0, 0, 1/5, 2/5, 3/5, 3/5, 4/5, 1, 1, 1\} \quad q = 2$ 







- Properties of surface basis functions
  - Extrema If p > 0 and q > 0,  $N_i^p(u) N_i^q(v)$  has a unique maximum.
  - Continuity Inside rectangles formed by the nodes  $u_i$  and  $v_i$ , the SF are infinitely differentiable. At a node  $u_i$  (resp.  $v_i$ ), $N_i^p(u) N_j^q(v)$  is (p-k) (resp. (q-k)) k)) times differentiable, k being the node multiplicity  $u_i$  (resp.  $v_i$ )
    - The continuity with respect to *u* (resp. *v*) depends solely on the nodal sequence U (resp. V).

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## Properties of surface basis functions

- A consequence of properties of the 1D shape functions
- Non-negativity
  - $N_{i}^{p}(u) N_{j}^{q}(v) \ge 0 \forall i, j, p, q, u, v$
- Partition of unity

$$\sum_{i=0}^{n} \sum_{j=0}^{m} N_{i}^{p}(u) N_{j}^{q}(v) = 1 \forall (u, v) \in [u_{min}, u_{max}] \times [v_{min}, v_{max}]$$

Compact support

$$\begin{split} N_i^p(u) N_j^q(v) &= 0 \text{ outside } (u, v) \in [u_i, u_{i+p+1}[\times[v_j, v_{j+q+1}[\\ \text{There are at most } (p+1)(q+1) \text{ non zero SF in a given} \\ \text{interval } [u_{i_0}, u_{i_0+1}[\times[v_{j_0}, v_{j_0+1}[\\ \dots \\ \text{In particular } N_i^p(u) N_j^q(v) \neq 0 \quad i_0 - p \leq i \leq i_0 \quad j_0 - q \leq j \leq j_0 \end{split}$$





- Properties of surface basis functions
  - Extrema If p>0 and q>0,  $N_i^p(u) N_i^q(v)$  has a unique maximum.
  - Continuity Inside rectangles formed by the nodes  $u_i$  and  $v_i$ , the SF are infinitely differentiable. At a node  $u_i$  (resp.  $v_j$ ),  $N_i^p(u) N_j^q(v)$  is (p-k) (resp. (q-k)) k)) times differentiable, k being the node multiplicity  $u_i$  (resp.  $v_i$ )
    - The continuity with respect to *u* (resp. *v*) depends solely on the nodal sequence U (resp. V).



 $U = \{0, 0, 0, 0, 1, 2, 2, 2, 2\} \quad p = 3 \quad U = \{0, 0, 0, 0, 1, 2, 3, 4, 5, 6, 6, 6, 6\}$  $V = \{0, 0, 0, 1, 2, 3, 3, 3\} \quad q = 2 \quad V = \{0, 0, 0, 1, 2, 3, 4, 5, 6, 7, 7, 7\}$ 



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- Computation of a point on the surface
  - 1 Find the nodal interval in which u is located

$$u \in [u_i, u_{i+1}]$$

2 – Compute the non vanishing 1D shape functions

$$N_{i-p}^{p}(u)$$
,  $\cdots$ ,  $N_{i}^{p}(u)$ 

- 3 Find the nodal interval in which v is located
- $v \in [v_j, v_{j+1}] \\ 4 \text{Compute the non vanishing 1D shape functions} \\ N_{j-q}^q(v), \cdots, N_j^q(v)$
- 5 Multiply the SFs with the adequate control points

$$S(u,v) = \sum_{k} \sum_{l} N_{k}^{p}(u) P_{kl} N_{l}^{q}(v) \quad i - p \le k \le i , j - q \le l \le j$$



 Each coloured square has an independent polynomial expression



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#### **B-Spline surfaces**



 $U = \{0, 0, 0, 0, 1, 2, 2, 2, 2\}$  p = 3 $V = \{0, 0, 0, 1, 2, 3, 3, 3\}$  q = 2





- Repetitions in the nodal sequence U or V
  - Discontinuities along iso-v or iso-u



 $U = \{0, 0, 0, 0, 2, 2, 2, 4, 4, 4, 4\} \qquad p = 3$  $V = \{0, 0, 0, 1.5, 1.5, 3, 3, 3\} \qquad q = 2$ 





- Properties of the B-Spline surface
  - Interpolate the 4 corners if the nodal sequences are of the form

$$U = \{\underbrace{0, \dots, 0}_{p+1}, u_{p+1}, \dots, u_{r-p-1}, \underbrace{1, \dots, 1}_{p+1}\}$$
$$V = \{\underbrace{0, \dots, 0}_{q+1}, v_{q+1}, \dots, v_{s-q-1}, \underbrace{1, \dots, 1}_{q+1}\}$$

If the nodal sequences correspond to Bézier curves  $: U = \{\underbrace{0, \cdots, 0}_{p+1}, \underbrace{1, \cdots, 1}_{p+1}\} \quad V = \{\underbrace{0, \cdots, 0}_{q+1}, \underbrace{1, \cdots, 1}_{q+1}\}$ 

then the surface is called a Bézier patch.





- Properties of the B-Spline surface
  - The surface has the property of affine invariance (invariance by translation in particular)
  - The convex hull of the control points contains the surface.
  - In every interval  $(u, v) \in [u_{i_0}, u_{i_0+1}] \times [v_{j_0}, v_{j_0+1}]$ , the portion of the surface is in the convex hull of the control points  $P_{ii}$ , (i, j) such that  $i_0 - p \le i \le i_0$   $j_0 - q \le j \le j_0$

- Control points may have a local control
- There is no variation diminishing property ( on the contrary to B-Spline/Bézier curves)





- Isoparametrics
  - Computation of isoparametrics is easy : Set u=u<sub>0</sub>

$$C_{u_{0}}(v) = S(u_{0}, v) = \sum_{i=0}^{n} \sum_{j=0}^{m} N_{i}^{p}(u_{0}) N_{j}^{q}(v) P_{ij}$$
  
$$= \sum_{j=0}^{m} N_{j}^{q}(v) \sum_{i=0}^{n} N_{i}^{p}(u_{0}) P_{ij} = \sum_{j=0}^{m} N_{j}^{q}(v) Q_{j}(u_{0})$$
  
with  $Q_{j}(u_{0}) = \sum_{i=0}^{n} N_{i}^{p}(u_{0}) P_{ij}$ 

• same with 
$$v = v_0$$
  
 $C_{v_0}(u) = S(u, v_0) = \sum_{i=0}^n N_i^p(u) Q_i(v_0)$   
with  $Q_i(v_0) = \sum_{j=0}^m N_j^q(v_0) P_{ij}$  18

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## **B-Spline surfaces**

- Derivatives of a B-Spline surface
  - We want to compute

$$\frac{\partial^{k+l}}{\partial u^k \partial v^l} S(u, v)$$

Differentiation of basis functions :

$$\frac{\partial^{k+l}}{\partial u^k \partial v^l} S(u, v) = \sum_{i=0}^n \sum_{j=0}^m \frac{\partial^{k+l}}{\partial u^k \partial v^l} N_i^p(u) N_j^q(v) P_{ij}$$

$$=\sum_{i=0}^{n}\sum_{j=0}^{m}\frac{\partial^{k}}{\partial u^{k}}N_{i}^{p}(u)\frac{\partial^{l}}{\partial v^{l}}N_{j}^{q}(v)P_{ij}$$

$$=\sum_{i=0}^{n}\sum_{j=0}^{m}N_{i}^{p(k)}(u)N_{j}^{q(l)}(v)P_{ij}$$





- Periodic surfaces
  - Like for the curves, possibility to "close" a B-Spline surface by transforming the nodal sequence
    - According to one parameter (u or v) Cylindrical surfaces
      - A single periodic nodal sequence
      - Control points on both sides of the seam are doubled
    - According to both parameters (u and v) Toroidal surfaces
      - Two periodic nodal sequences
      - Some control points are repeated 4 times !







#### Möbius strip







 $U = \{-3, -2, -1, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15\}$  p=3 $V = \{0, 0, 0, 1, 2, 3, 4, 5, 6, 7, 7, 7\}$  p = 222







 $U = \{-3, -2, -1, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15\}$  p=3 $V = \{0, 0, 0, 1, 2, 3, 4, 5, 6, 7, 7, 7\}$  p = 223





















 $U = \{-3, -2, -1, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15\}$  p=3 $V = \{-2, -1, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$  p = 227







- Some operations
  - Insertion of nodes
    - Extraction of iso-parametrics
    - Obtaining the position of a point on the surface
    - Subdivision of the surface
    - Transformation into Bézier patches

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#### Insertion of nodes

- We insert nodes In a nodal sequence (U ou V)
- The new nodal sequence replace the old one
- The control points are modified
  - If U is modified, every series of control points corresponding to v=cst is independently modified
  - If V is modified, every series of control points corresponding to u=cst is independently modified
  - We use Boehm's algorithm as for curves

$$V = \{0, 0, 0, 1, 2, 3, 3, 3\} p = 2$$

$$V = \{0, 0, 0, 0, 1, 2, 3, 3, 3, 3\} p = 3$$

$$U = \{0, 0, 0, 0, 1, 2, 3, 3, 3, 3\} p = 3$$













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- Extraction of iso-parametrics using node insertion
  - We must saturate one node in  $u=u_{iso}$  (resp. in  $v=v_{iso}$ ).
  - The new control points obtained by Boehm's algorithm do form the control polygon of the iso-parametric curve.
  - The nodal sequence of this curve is V (resp. U).













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Obtaining the position of a point on the surface



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### Subdivision into Bézier patches

... one must saturate every node of each nodal sequence U and V



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### Continuity requirements for surfaces

 C<sup>1</sup> vs C<sup>2</sup> – becomes visible when light interaction comes into play







#### **CAD** Surfaces





- Parametric surfaces: an explicit representation
  - Lightweight
  - Discretization algorithms are non trivial... but it is necessary for display purposes and in computer graphics
  - Generally, these surfaces are used in cases where the geometric accuracy is essential, as in the computation of intersections and other precise geometric primitives

### Modeling operators are not trivial

 In computer graphics, such accuracy is generally not needed.





- Subdivision surfaces
  - Modelling basis = elementary mesh
  - By successive iterations, this mesh is refined up to the accuracy needed for the application
  - It is more like an algorithmic description vs. an algebraic representation, because the algorithm that is used to subdivide the mesh determines the final shape and the properties of the limiting surface (*ie.* when the number of subdivisions tends to the infinite)
  - Some of these limiting surfaces are equivalent to "regular" parametric surfaces, therefore have the same "accuracy".





### History

1974 – George Chaikin

An algorithm for high speed curve generation

1978 – Daniel Doo & Malcolm Sabin

(D) A subdivision algorithm for smoothing irregularly shaped polyhedrons

(D&S) Behaviour of recursive division surfaces near extraordinary points.

1978 – Edwin Catmull & Jim Clark

Recursively generated B-Spline surfaces on arbitrary topological meshes

1987 – Charles Loop

Smooth subdivision surfaces based on triangles

2000 – Leif Kobbelt

 $\sqrt{3}$  – subdivision (interpolating scheme)





Chaikin's scheme



George Chaikin, An algorithm for high speed curve generation, *Computer graphics and Image Processing* **3** (1974), 346-349





Chaikin's scheme

or « Corner-cutting »







Chaikin's scheme



Chaikin's idea was simple : repeating the corner cutting, to the limit, one obtains a smooth curve

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- Chaikin's scheme
  - Starting from a polygon having *n* vertices  $\{P_0, P_1, \dots, P_{n-1}\}$ , one builds the polygon having 2n vertices  $\{Q_0, R_0, Q_1, R_1, \dots\}$  $Q_{n-1}, R_{n-1}$ . This polygon serves as a basis for the next step of the algorithm :  $\{P'_{0}, P'_{1}, \dots, P'_{2n-1}\}$  $Q_i = \frac{3}{4}P_i + \frac{1}{4}P_{i+1}$







- Chaikin's scheme
  - Riesenfeld (1978) has shown that this algorithm leads at the limit to an uniform quadratic B-spline, which exhibits a  $C^1$  continuity.







- Demonstration of the equivalence of Chaikin's scheme and uniform quadratic B-Splines
  - The B-Spline curve is defined by :  $P(u) = \sum_{i=0}^{n} P_i N_i^2(u)$

$$N_{i}^{2}(u) = \frac{u - u_{i}}{u_{i+2} - u_{i}} N_{i}^{1}(u) + \frac{u_{i+3} - u_{i}}{u_{i+3} - u_{i+1}} N_{i+1}^{1}(u)$$

$$N_{i}^{1}(u) = \frac{u - u_{i}}{u_{i+1} - u_{i}} N_{i}^{0}(u) + \frac{u_{i+2} - u}{u_{i+2} - u_{i+1}} N_{i+1}^{0}(u)$$
$$N_{i}^{0}(u) = \begin{cases} 1 & \text{if} \\ 0 & \text{otherwise} \end{cases} \quad u_{i} \le u < u_{i+1} \end{cases}$$

$$U = \{u_0, \dots, u_{n+2}\}, u_{i+1} - u_i = 1, i = 0 \cdots n + 1$$

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It can be rewritten as an assembly of curved "parts" :

$$P(u) = \sum_{i=0}^{n} P_i N_i^2(u) = \sum_{k=0}^{n-2} P_k(u)$$
 « Portion » of curve

with 
$$P_k(u) = \begin{bmatrix} 1 & u & u^2 \end{bmatrix} \cdot M_k \cdot \begin{bmatrix} P_k \\ P_{k+1} \\ P_{k+2} \end{bmatrix}$$

• The matrix  $M_k$  depends on the nodal sequence U.













- Binary subdivision of a B-Spline curve for  $0 \le u \le 1$ 
  - One has to find the new set of control points for each halves of the curve

We set n=2 (number of control points -1)

$$P(u) = \begin{bmatrix} 1 & u & u^2 \end{bmatrix} \cdot M \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix} \qquad \qquad M = \frac{1}{2} \begin{bmatrix} 1 & 1 & 0 \\ -2 & 2 & 0 \\ 1 & -2 & 1 \end{bmatrix}$$

• One wants to express  $P_{[0,1/2]}(u)$  and  $P_{[1/2,1]}(u)$ - on each subdivision, the parameter u shall be between 0 and 1.  $P_0$   $P_{[0,1/2]}(u)$  $P_{[1/2,1]}(u)$ 



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• Case of  $P_{[0,1/2]}(u)$ 

$$P_{[0,1/2]}(u) = P(u/2) = \begin{bmatrix} 1 & u/2 & u^2/4 \end{bmatrix} \cdot M \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & u & u^2 \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1/2 & 0 \\ 0 & 0 & 1/4 \end{bmatrix} \cdot M \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & u & u^2 \end{bmatrix} \cdot M \cdot M^{-1} \cdot \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1/2 & 0 \\ 0 & 0 & 1/4 \end{bmatrix} \cdot M \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & u & u^2 \end{bmatrix} \cdot M \cdot \begin{bmatrix} Q_0 \\ Q_1 \\ Q_2 \end{bmatrix} \text{ avec } \begin{bmatrix} Q_0 \\ Q_1 \\ Q_2 \end{bmatrix} = \underbrace{M^{-1} \cdot \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1/2 & 0 \\ 0 & 0 & 1/4 \end{bmatrix} \cdot M \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix}$$



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• Case of  $P_{[1/2,1]}(u)$ 

$$P_{[1/2,1]}(u) = P((1+u)/2) = \begin{bmatrix} 1 & (1+u)/2 & (1+u)^2/4 \end{bmatrix} \cdot M \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix}$$
$$= \begin{bmatrix} 1 & u & u^2 \end{bmatrix} \cdot \begin{bmatrix} 1 & 1/2 & 1/4 \\ 0 & 1/2 & 1/2 \\ 0 & 0 & 1/4 \end{bmatrix} \cdot M \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix}$$
$$= \begin{bmatrix} 1 & u & u^2 \end{bmatrix} \cdot M \cdot M^{-1} \cdot \begin{bmatrix} 1 & 1/2 & 1/4 \\ 0 & 1/2 & 1/2 \\ 0 & 0 & 1/4 \end{bmatrix} \cdot M \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix}$$
$$= \begin{bmatrix} 1 & u & u^2 \end{bmatrix} \cdot M \cdot \begin{bmatrix} R_0 \\ R_1 \\ R_2 \end{bmatrix} \quad \text{with} \quad \begin{bmatrix} R_0 \\ R_1 \\ R_2 \end{bmatrix} = \underbrace{M^{-1} \cdot \begin{bmatrix} 1 & 1/2 & 1/4 \\ 0 & 1/2 & 1/2 \\ 0 & 0 & 1/4 \end{bmatrix} \cdot M \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix}$$





#### • Finally,

$$\begin{bmatrix} Q_0 \\ Q_1 \\ Q_2 \end{bmatrix} = S_{[0,1/2]} \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix} \qquad S_{[0,1/2]} = M^{-1} \cdot \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1/2 & 0 \\ 0 & 0 & 1/4 \end{bmatrix} \cdot M = \frac{1}{4} \begin{bmatrix} 3 & 1 & 0 \\ 1 & 3 & 0 \\ 0 & 3 & 1 \end{bmatrix} \qquad \begin{bmatrix} Q_0 \\ Q_1 \\ Q_2 \end{bmatrix} = \frac{1}{4} \begin{bmatrix} 3 P_0 + P_1 \\ P_0 + 3 P_1 \\ 3 P_1 + P_2 \end{bmatrix}$$
$$\begin{bmatrix} R_0 \\ R_1 \\ R_2 \end{bmatrix} = S_{[1/2,1]} \cdot \begin{bmatrix} P_0 \\ P_1 \\ P_2 \end{bmatrix} \qquad S_{[1/2,1]} = M^{-1} \cdot \begin{bmatrix} 1 & 1/2 & 1/4 \\ 0 & 1/2 & 1/2 \\ 0 & 0 & 1/4 \end{bmatrix} \cdot M = \frac{1}{4} \begin{bmatrix} 1 & 3 & 0 \\ 0 & 3 & 1 \\ 0 & 1 & 3 \end{bmatrix} \qquad \begin{bmatrix} R_0 \\ R_1 \\ R_2 \end{bmatrix} = \frac{1}{4} \begin{bmatrix} P_0 + 3 P_1 \\ 3 P_1 + P_2 \\ P_1 + 3 P_2 \end{bmatrix}$$



One finds the same coefficients as in Chaikin's scheme...

(except for the indices)

$$Q_{i} = \frac{3}{4} P_{i} + \frac{1}{4} P_{i+1}$$
$$R_{i} = \frac{1}{4} P_{i} + \frac{3}{4} P_{i+1}$$

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- This can be extended to cubic B-Splines
  - $C^2$  continuity









- Doo-Sabin scheme
  - This is an extension of Chaikin's scheme for a uniform biquadratic B-Spline surface
  - The new mesh is built using the control points resulting from the subdivision of the original patch into 4 new subpatches.







 Expression of the bi-quadratic patch as a monomial form for 0 ≤ u ≤ 1 and 0 ≤ v ≤ 1:





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• Subdivision - patch  $S_1(u,v)$ 

$$S_{1}(u,v) = S(u/2,v/2) = \begin{bmatrix} 1 & u/2 & u^{2}/4 \end{bmatrix} \cdot M \cdot P \cdot M^{\mathrm{T}} \cdot \begin{bmatrix} 1 \\ v/2 \\ v^{2}/4 \end{bmatrix} \qquad P = \begin{bmatrix} P_{00} & P_{01} & P_{02} \\ P_{10} & P_{11} & P_{12} \\ P_{20} & P_{21} & P_{22} \end{bmatrix}$$
$$S_{1}(u,v) = S(u/2,v/2) = \begin{bmatrix} 1 & u & u^{2} \end{bmatrix} \cdot C \cdot M \cdot P \cdot M^{\mathrm{T}} \cdot C^{\mathrm{T}} \begin{bmatrix} 1 \\ v \\ v^{2} \end{bmatrix} \qquad C = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1/2 & 0 \\ 0 & 0 & 1/4 \end{bmatrix}$$
$$= \begin{bmatrix} 1 & u & u^{2} \end{bmatrix} \cdot M \cdot M^{-1} \cdot C \cdot M \cdot P \cdot M^{\mathrm{T}} \cdot C^{\mathrm{T}} \cdot (M^{-1})^{\mathrm{T}} \cdot M^{\mathrm{T}} \cdot \begin{bmatrix} 1 \\ v \\ v^{2} \end{bmatrix}$$

$$P_{00} = \begin{bmatrix} 1 & u & u^{2} \end{bmatrix} \cdot M \cdot (M^{-1} \cdot C \cdot M) \cdot P \cdot (M^{-1} \cdot C \cdot M)^{\mathrm{T}} \cdot M^{\mathrm{T}} \cdot \begin{bmatrix} 1 \\ v \\ v^{2} \end{bmatrix}$$
$$= \begin{bmatrix} 1 & u & u^{2} \end{bmatrix} \cdot M \cdot P' \cdot M^{\mathrm{T}} \cdot \begin{bmatrix} 1 \\ v \\ v^{2} \end{bmatrix} P' = S \cdot P \cdot S^{\mathrm{T}} \qquad S = M^{-1} \cdot C \cdot M$$

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$$P' = \frac{1}{16} \begin{bmatrix} 3(3P_{00} + P_{10}) + 3P_{01} + P_{11} & 3P_{00} + P_{10} + 3(3P_{01} + P_{11}) & 3(3P_{01} + P_{11}) + 3P_{02} + P_{12} \\ 3(P_{00} + 3P_{10}) + P_{01} + 3P_{11} & P_{00} + 3P_{10} + 3(P_{01} + 3P_{11}) & 3(P_{01} + 3P_{11}) + P_{02} + 3P_{12} \\ 3(3P_{10} + P_{20}) + 3P_{11} + P_{21} & 3P_{10} + P_{20} + 3(3P_{11} + P_{21}) & 3(3P_{11} + P_{21}) + 3P_{12} + P_{22} \end{bmatrix}$$

 Same developments should be done with the 3 other quadrants, and will lead to the same "structure"



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### LIÈGE CAD & Computational Geometry Subdivision surfaces

• Subdivision - patch  $S_2(u,v)$ 

$$\begin{split} S_{2}(u,v) &= S(u/2,(1+v)/2) = [1 \ u/2 \ u^{2}/4] \cdot M \cdot P \cdot M^{\mathrm{T}} \cdot \begin{bmatrix} 1\\(1+v)/2\\(1+v)^{2}/4 \end{bmatrix} \quad P = \begin{bmatrix} P_{00} \ P_{01} \ P_{02}\\P_{10} \ P_{11} \ P_{12}\\P_{20} \ P_{21} \ P_{22} \end{bmatrix} \\ S_{2}(u,v) &= S(u/2,(1+v)/2) = [1 \ u \ u^{2}] \cdot C_{u} \cdot M \cdot P \cdot M^{\mathrm{T}} \cdot C_{v}^{\mathrm{T}} \begin{bmatrix} 1\\v\\v^{2} \end{bmatrix} \\ C_{u} &= \begin{bmatrix} 1 \ 0 \ 0\\0 \ 1/2 \ 0\\0 \ 0 \ 1/4 \end{bmatrix} \quad C_{v} = \begin{bmatrix} 1 \ 1/2 \ 1/4\\0 \ 1/2 \ 1/2\\0 \ 0 \ 1/4 \end{bmatrix} \\ &= [1 \ u \ u^{2}] \cdot M \cdot (M^{-1} \cdot C_{u} \cdot M) \cdot P \cdot (M^{-1} \cdot C_{v} \cdot M)^{\mathrm{T}} \cdot M^{\mathrm{T}} \cdot \begin{bmatrix} 1\\v\\v^{2} \end{bmatrix} \\ &= [1 \ u \ u^{2}] \cdot M \cdot Q' \cdot M^{\mathrm{T}} \cdot \begin{bmatrix} 1\\v\\v^{2} \end{bmatrix} \quad Q' = S_{u} \cdot P \cdot S_{v}^{\mathrm{T}} \quad S_{u} = M^{-1} \cdot C_{u} \cdot M \\ S_{v} = M^{-1} \cdot C_{v} \cdot M \end{split}$$



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#### **Subdivision surfaces**

$$Q' = S_{u} \cdot P S_{v}^{T} \qquad S_{u} = M^{-1} \cdot C_{u} \cdot M = \frac{1}{4} \begin{bmatrix} 3 & 1 & 0 \\ 1 & 3 & 0 \\ 0 & 3 & 1 \end{bmatrix} \qquad P_{00} \qquad \qquad P_{02} \qquad \qquad P_{02} \qquad \qquad P_{02} \qquad \qquad P_{01} \qquad \qquad P_{02} \qquad \qquad P_{0$$

Some of the points (2) are already computed, e.g.

$$Q'_{00} = 3(P_{00} + 3P_{01}) + P_{10} + 3P_{11}$$
  
( =  $P'_{01} = 3P_{00} + P_{10} + 3(3P_{01} + P_{11})$  )





• Extension to meshes showing an arbitrary topology P

topology  $P_{00}$   $P_{02}$   $P_{00}$   $P_{02}$   $P_{02}$  $P_{02}$ 

- The new vertices are obtained as a simple arithmetic mean of 3 categories of vertices :
  - The vertices of the old mesh
  - Vertices on the edges (barycenter of the extremities of the edge)
  - Vertices inside a face (barycenter of the vertices of the face)





### Extension to meshes showing an arbitrary topology

1 – Computation of the vertices P' (for each vertex P, compute the mean between P, the vertices on the adjacent faces, and the vertices on adjacent edges)







- Extension to meshes showing an arbitrary topology
  - 2 For each face, link the corresponding vertices P'







#### Extension to meshes showing an arbitrary topology

3 – For each old vertex, connect the new ones that have been created for each adjacent face to this old vertex.







### Extension to meshes showing an arbitrary topology

4 – For each old edge, connect the new vertices that have been created for each adjacent face to this old edge.







- Some points on the mesh and the limiting surface are « extraordinary »
  - These are vertices with a valence (number of incident edges) that is different from 4.



• Everywhere the continuity of the limiting surface is  $C^1$ ; except at extraordinary points, where it decreases to  $C^0$ .





- Catmull-Clark scheme
  - Similar idea for bicubic B-Splines (proof : Jos Stam, Siggraph 1998)  $P_{nn}$







### Three types of vertices

- « Face » vertices are at the barycentre of the vertices of that face:

 $P_f = Q$ 

- « Edge » vertices are at the barycentre of the extremities of the edge and the two « Face » vertices of the adjacent faces :

$$P_e = \frac{Q+R}{2}$$

- « Corner » vertices are positioned such that :

 $P_v = \frac{Q + 2R + (n-3)S}{n}$ 

Q = mean of the barycentre of the incident faces R = mean of the barycentre of the incident edges S = original vertex

*n* = number of incident edges to *S* (Valence)



- « face » vertices
- « edge » vertices
- « corner » vertices

#### LIEGE CAD & Computational Geometry **Subdivision surfaces** 1/4 1/46/64 1/64 1/100 6/100 36/64 1/41/46/64 $\overline{\bigcirc}$ Face 6/100 65/100



vertex



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Reconnecting the new vertices

1 – Connect the « face » vertices to the « edge » vertices of neighbouring edges

2 – Connect the « corner » vertices to the « edge » vertices of incident edges






#### **Subdivision surfaces**

• As with Doo-Sabin surface, the continuity is degraded for some extraordinary vertices. The bicubic surfaces are therefore  $C^2$  everywhere except at extraordinary points : it is then only  $C^1$ .







#### **Subdivision surfaces**

- Loop's scheme
  - Allows to subdivide triangular meshes
  - The limiting surface is  $C^2$ , except at extraordinary vertices of valence <>6, where it is only  $C^1$ .
  - The principle is to subdivide triangles into 4 sub-triangles.
    - Corner vertices (black) and edge vertices (red) are created.







- As such, works only for vertices with a valence equal to 6
- It may be extended to other valences, but the formula has to be adapted such that the resulting surface is smooth.

• Let 
$$V^1 = \alpha_n V + (1 - \alpha_n) Q$$
  
with  $\alpha_n = \left(\frac{3}{8} + \frac{1}{4} \cos \frac{3\pi}{n}\right)^2 + \frac{3}{8}$ 

n is the valence of the original vertex.

 On boundaries : vertices should not move inside the surface, they should rather slide along the boundary. One recovers the classical cubic B-Spline scheme in that case







#### **Subdivision surfaces**

Loop's scheme







#### **Subdivision surfaces**

- Kobbelt's  $\sqrt{3}$  scheme
  - See paper on the course's website
  - For a similar level of refinement, it generates less triangles than Loop's scheme







#### Subdivision surfaces in CATIA

- A very easy-touse design tool
- As S-s are equivalent to some class of B-Spline surfaces, they retain a good degree of accuracy
   "CATIA Shape" module Imagine Shape (IMA) tool



#### LIÈGE CAD & Computational Geometry Subdivision surfaces



- ... and in NX
- A very easy-touse design tool
- As S-s are equivalent to some class of B-Spline surfaces, they retain a good degree of accuracy
   "NX Realize
- shape" tool



#### LIÈGE CAD & Computational Geometry Université Splines of all kinds



- Zoo of Splines ...
- Lg-splines Analytic splines Parabolic Arc splines Beta splines **B**-splines Bernoulli splines Box splines Cardinal splines Circular splines Complete splines Nu-splines Natural splines L-splines Whittaker splines

Nonlinear splines D<sup>m</sup> splines Discrete splines Euler splines Exponential splines Gamma splines **GB-splines HB-splines** Hyperbolic Splines **Complete Monosplines** Tchebycheffian splines **Tension splines** Trigonometric splines Bézier splines

One-sided splines Parabolic splines Perfect splines Periodic splines Poly-splines Rational splines Simplex splines Spherical splines Taut splines Complex splines Confined splines Deficient splines Thin plate splines





#### Solid models and B-REP

## LIÈGE CAD & Computational Geometry Solid modelling



#### Classical modelling problem : the intersection



3 independent representations of the intersection :

- a 3D NURBS curve (giving points in the global XYZ coordinate system)

- a 2D NURBS curve

in the parametric space of surface A (giving 2D points in the coordinate system of the parametric space of surface A)

- Idem for surface B



Parametric space of surface B



#### Solid modelling

- Theoretically, these three representations are equivalent ... 3 independent representations of the intersection :
  - a 3D NURBS curve
    - a 2D NURBS curve (parametric space of surface A)
    - a 2D NURBS curve (parametric space of surface B)
- In practice, there are numerical approximations
  - NURBS are finite approximation spaces; therefore approximation/interpolation errors do occur.
  - The use of floating point numbers with a finite binary representation of the mantissa lead to numerical errors
- There is no robust way to ensure, in a geometrical sense, that a curve located on surface A is the same as the corresponding curve on surface B, *i.e.* that both surfaces are neighbours, and share the same edge.

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- Definition of a topology : non geometric relations between entities.
- This allows to unify the calculations (of points, normals, etc...) on entities shared (or bounding) other entities (*eg.* an edge shared by surfaces).
- It also allows the explicit definition of volumes from the surfaces that bound the volume.
- It may also solve the problem of orientation of surfaces





#### Solid modelling

Topology

#### B-REP model



## **LIEGE CAD & Computational Geometry**



#### Solid modelling

- B-Rep model
  - « Boundary representation »
  - Model based on the representation of surfaces
  - Model of exchange (STEP format) and definition
  - The "natural" set of operators is richer than for CSG
    - Extrusion, chamfer etc ...
  - Does not carry the history of construction of the model (whereas CSG usually does)



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#### B-Rep model

- Consists of two types of information :
  - Geometric
    - Geometric information is used for defining the spatial position, the curvatures, etc...
    - That's what we have seen until now NURBS curves and surfaces !
  - Topological

This allows to make links between geometrical entities.

- Two types of entities
  - Geometric entities: (volume), surface, curve, point
  - Topological entities : solid, face, edge, vertex
- A topological entity "lies on" a geometric entity, which is its geometrical support (when existing)

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- B-Rep model
  - Complete hierarchical model







#### Solid modelling



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#### LIÈGE CAD & Computational Geometry Solid modelling









#### Solid modeling



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How to obtain the (x,y,z) coordinates of the encircled point?

1 - Use the 3D vertex directly (Point xyz xxx)

2 – Use the boundary vertices for every 3D edge (there are 3 such edges ) (Point\_t t1,t2,t3) Then use those (t) to get (x,y,z) by the 3D edges

3 – Use the boundary vertices of the 2D boundary edges in the face (there are 2 faces, so 4 of them), (Point t t'1, t'2, t'3, t'4). Then use those (t) to obtain coordinates (u,v) in the parametric space of the face, thanks to 2D curves, finally, use those (u,v) to obtain (x,y,z) thanks to the geometry of the face.

So there exists 8 different ways. Nothing indicates that the 8 set of 3D coordinates are exactly equal (there are numerical approximations). Only topology allows us to say that those 8 points are all referring to the same point ... at least conceptually.







#### Solid modeling

- B-Rep model
  - Euler characteristic for polyhedra

 $\chi(S) = v - e + f$ 

Euler – Poincaré formula  $\chi(S) = v - e + f - r = 2(s - h)$ 

with

r

- = number of vertices
- of faces
- ... of edges *e* =
- 11 of solids (independent volumes) *s* =
- h =... of holes – going through (topol. gender)
  - of internal loops (rings) 11





#### **Euler's formula**

#### Euler characteristic





Example : Cube  $v - e + f = \kappa$ 

Opened and flattened Cube  $v - e + f = \kappa -1$ 

Step 0 : we take a face off the polyhedron and flatten it to obtain a plane graph





#### **Euler's formula**



 $v - e + f = \kappa - 1$  +1e, +1f +5e, +5f  $v - e + f = \kappa - 1$   $v - e + f = \kappa - 1$ 

Step 1 : Repeat the following operation :

For each non triangular face, add one edge linking non related vertices. Each tome, the number of edges and faces is increased by 1. This is repeated until no non triangular faces remain.





#### **Euler's formula**



 $v - e + f = \kappa - 1$  -1e, -1f -2e, -1f, -1v $v - e + f = \kappa - 1$   $v - e + f = \kappa - 1$ 

Step 2 : One alternates between these two operationsPreferentially, delete triangles that have 2 boundary edges.Every time; *e* decreases by 2 and *f* and *v* by 1.

- Then, delete triangles with only one boundary edege. Each tile, *e* and *f* decrease by 1.

This until only one triangle remain.



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$$v - e + f = \kappa - 1$$

#### LIÈGE CAD & Computational Geometry Euler's formula



# Every polygon can be decomposed into triangles Therefore, by applying the three operations described in the previous slides, we can transform the planar graph into a triangle without changing Euler's characteristic. The triangle obviously satisfies

$$v - e + f = \kappa - 1$$

with  $\kappa$ -1 = 1

As a consequence, the planar graph verifies the formula.

• So the initial polyhedron satisfies :

$$v - e + f = \kappa = 2$$



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Necessity to take "rings" into account - inside faces



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Necessity to take "holes" into account









polyhedron  $\chi(S)=2$   $\chi(S)=v-e+f$  $\chi(S)=8-12+6=2$ 

"Warped" polyhedron  $\chi(S)=8-12+6=2$ 

OK

Polyhedron with one hole, 4 edges less, 2 faces k=2 less, 4 vertices less  $\chi(S)=4-8+4\neq 2$ Not OK ! Contribution of the hole  $\chi(S)=v-e+f-r+2h=2$ 

#### LIÈGE CAD & Computational Geometry Université Solid modeling



 Every B-rep model is identifiable (topologically) to a « point » in a 6-dimensional vector space.

Vector space of coordinates v, e, f, s, h, r.

- Any topologically valid model shall verify the Euler-Poincaré relation
  - This relation defines an « hyperplane » ( of dimension 5) in a 6-dimensional space
  - The equation of this hyperplane is :

v - e + f - 2s + 2h - r = 0





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$$v - e + f - 2s + 2h - r = 0$$

We can update a valid solid and modify the 6 numbers characterising a model with a transformation that yields a valid solid for which :

$$v + \Delta v - e - \Delta e + f + \Delta f$$
  
-2s-2\Delta s+2h+2\Delta h-r-\Delta r=0  
$$\Rightarrow \Delta v - \Delta e + \Delta f - 2\Delta s + 2\Delta h - \Delta r = 0$$

- In this way, add a vertex ( $\Delta v = 1$ ) must be accompanied, one way or another, by addition of an edge ( $\Delta e = 1$ ) OR of the withdrawal of a face (  $\Delta f = -1$ ), etc...
- Elementary operations satisfying the Euler-Poincaré relation are called **Euler operators**.
- They allow staying on the « hyperplane » of validity while changing the topological configuration

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#### Solid modelling

- Euler operators
  - The use of Euler operators guarantees the topological validity of the result
  - Here we don't check the geometric validity (selfintersections etc...)
  - We identify them under the form : MaKb where M = Make **K** = Kill and **a** and **b** are a sequence of entities : vertex, edge, face, solid, hole or ring.
    - In total, there are 99 Euler operators aiming to modify the number of entities by at most one unit.
      - These are divided in 49 + 49 inverses, plus the identity operator.
    - Among those 49 operators, we can chose 5 linearly independent operators (the hyperplane has 5 dimensions)
    - Those 5 independent operators form a base for the hyperplane of topologically admissible models

## LIEGE CAD & Computational Geometry Solid modelling



#### Example of a set of Euler operators



MEF, Make an Edge and a Face MEKR, Make an Edge and Kill a Ring MVFS, Make a Vextex, a Face and a Shell KFMRH, Kill a Face Make a Ring and a Hole **KEV**, Kill an Edge and a Vertex KEF, Kill an Edge and a Face KEMR, Kill an Edge and Make a Ring KVFS, Kill a Vertex, a Face and a Shell

- MFKRH, Make a Face, Kill a Ring and a Hole
- Proof by Mäntylä (1984) that those operators allow to build every valid solid (since they are independent)
- Those operators form a base of the space of valid configurations (the « hyperplane »)

#### LIÈGE CAD & Computational Geometry Solid modelling



There are three types of operators in this set :

- Skeleton operators MVFS and KVFS
  - Allow to build/destroy elementary volumes
- Local operators MEV, KEV, MEF, KEF, KEMR, MEKR
  - Allow to modify connectivities for existing volumes
  - Don't modify fundamental topological characteristics of the surfaces - nb of handles/ holes (topological gender) and number of independent volumes
- Global operators KFMRH and MFKRH
  - Allow to add / remove "handles" (change the topological gender)
- Only the skeleton and global operators do change the topological gender.

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#### Solid modelling

- Euler operators
  - « Skeleton » operators

#### MVFS; KVFS

Allow to « build » an « elementary » volume from void (which is an admissible topological structure) – or destroy it.







#### Solid modelling

- Euler operators
  - Local operators
     MEV, KEV (case 1)






- Euler operators
  - Local operators
    MEV, KEV (case 2)







- Euler operators
  - Local operators
    MEV, KEV (case 3)







- Euler operators
  - Local operators
    MEF, KEF (case 1)







- Euler operators
  - Local operators
    MEF, KEF (case 2)







- Euler operators
  - Local operators
    MEF, KEF (case 3)







- Euler operators
  - Local operators
    KEMR, MEKR (case 1)







- Euler operators
  - Local operators

**KEMR**, **MEKR** (case 2 : the loop for the internal ring is reduced to a single vertex )







- Euler operators
  - Local operators

**KEMR**, **MEKR** (case 3 : both loops are reduced to one vertex – one is the ring; the other is the external loop of the face)







- Euler operators
  - Global operators

KFMRH, MFKRH (case 1 : allow the creation / destruction of holes in a solid)







- Euler operators
  - Global operators

**KFMRH**, **MFKRH** (case 2) : join two independent solids : here more judiciously called Kill Face, Solid and Make Ring (KFSMR)

Interpretation of global operators is sometimes confusing







# Example of use of Euler operators







## Example of use of Euler operators







## Example of use of Euler operators



S. Havemann<sup>121</sup>

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## Solid modelling

- Those operators have a vectorial form in the basis of elementary entities
  - vefhrs
  - (1, 1, 0, 0, 0, 0) MEV, Make an Edge and a Vertex
  - (0, 1, 1, 0, 0, 0) MEF, Make a Face and an Edge
  - (0,-1,0,0,1,0) KEMR, Kill an Edge Make a Ring
  - (1, 0, 1, 0, 0, 1) MVFS, Make a Vertex, a Face and a Solid
  - (0, 0, -1, 1, 1, 0) KFMRH, Kill a Face, Make a Ring and a Hole
  - In order to have a complete basis of the configuration space, a vector orthogonal to the hyperplane of acceptable configurations must be added

$$v - e + f - 2s + 2h - r = 0$$

The coefficients of the equation of hyperplane are precisely the coordinates of the orthogonal vector... v e f h r s

```
(1,-1, 1, 2,-1,-2) – Euler-Poincaré
```

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- Any transformation can thus be expressed easily using matrix operations
  - A is a basis of the topological configurations space
  - The columns of A are the variation of the number of entities for each operator, and the E-P relation.







#### A is composed of linearly independent vectors, thus one can get the inverse...



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Determination of elementary operations ...







Are Euler coordinates sufficient to define the topology of a solid  $? \rightarrow No$ .



1 x MEF, Make an Edge and a Face 1 x KEF, Kill an Edge and a Face

7 x MEV, Make an Edge and a Vertex 7 x MEF, Make an Edge and a Face 1 x MVFS, Make a Vertex, a Face and a Solid 0 x KEMR, Kill an Edge Make a Ring 0 x KFMRH, Kill a Face, Make a Ring and a Hole

7 x MEV 7 x MEF 1 x MVFS 0 x KEMR \_0 x KFMRH

Identical Euler coordinates





### Each Euler operator takes a certain number of parameters, in principle the entities to destroy/ or replace, and data necessary to creation of new entities.

- These depend on the structure of data used to represent the B-Rep object
- Application of an Euler operator is not always possible, the entities involved must exist and respect some conditions

KEF for example may only be applied on an edge separating two distinct faces... if not, one does not remove any face from the model !

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### Conditions of application of Euler operators

MEV, Make an Edge and a Vertex	No specific conditions except existence of a solid
KEV, Kill an Edge and a Vertex	The edge has two distinct vertices
MEF, Make an Edge and an Face	Vertices belonging to the same boundary loop of one face
KEF, Kill an Edge and a Face	Distinct faces located on both sides of the edge
MEKR, Make an Edge and Kill a Ring	Vertices belong to distinct boundary loops of the same face
KEMR, Kill an Edge and Make a Ring	Same face located on both sides of the edge, which is not part of a ring
MVFS, Make a Vextex, a Face and a Shell	Empty space
KVFS, Kill a Vertex, a Face and a Shell	The shell (solid) has no edges and has only one vertex (elementary volume)
KFMRH, Kill Face Make a Ring and a Hole	The face cannot hold any ring
MFKRH, Make a Face, Kill a Ring and a Hole	May be only applied to a ring $128$





- Some examples of the application of Euler operators (not shown here)
  - Extrusion of a face
  - Junction of two solids
  - Cutting out a solid by a plane
  - Boolean operations between solids





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